

Today

- 1) 2 common setups
- 2) Concentration Framework
- 3) Solving (1) w/ (2)

Admin

- Scribe Note Sign-up
- Hw 1 in
- Hw 2 out
- Caddy
- Names

2 Scenarios

Balls $\rightarrow$ bins uniformly at random + independently

o o o o ... o n balls

□ □ □ ... □ m bins

1) If $m=n$, what is max # balls in a bin?

$$\sim \ln n \quad (\text{w/ pr} \geq 1 - \frac{1}{n})$$

2) How big must n be so ≥ 1 ball in each bin?

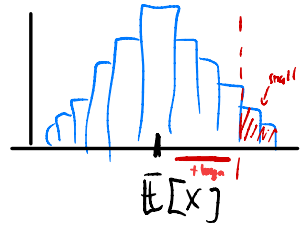
$$\sim n \ln n \quad (\text{for pr} \geq 1 - \frac{1}{n})$$

"coupon collector"

How to Solve Your Fave Randomization Question

Concentration Framework

- a) Show (*) true if all RVs near $\mathbb{E}$
- b) Concentration: one RV at $\mathbb{E} (\pm \log n)$ w/ high pr
- c) Union Bound: all RVs at $\mathbb{E} (\pm \log n)$ w/ high pr
↳ Reminder: $\Pr(B_1 \cup B_2 \cup \dots) \leq \sum_i \Pr(B_i)$



Let $X_{ij} :=$ indicator of ball $j \rightsquigarrow$ bin i

Let $X_i := \sum_j X_{ij}$ (# balls in bin i)

1) $\mathbb{E}[X_i] = 1$ (when $m = n$)

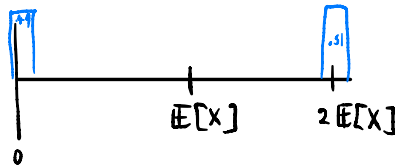
2) $\mathbb{E}[X_i] = \Omega(\log n)$ (when $m \geq \Omega(n \cdot \log n)$)

Markov's Inequality

Suppose X is a non-negative RV w/ $\mathbb{E}[X] > 0$

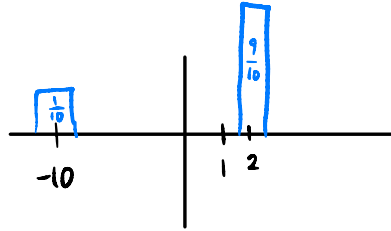
$$\Pr(X \geq a \cdot \mathbb{E}[X]) \leq \frac{1}{a} \rightarrow \text{at least } \frac{1}{a} \text{ of the time}$$

Intuition



Why Non-Negative?

$$\mathbb{E}[X] = \frac{4}{5}$$



$$a \cdot \mathbb{E}[X] \cdot \Pr(X \geq a \cdot \mathbb{E}[X]) \leq a \cdot \mathbb{E}[X] \cdot \sum_{i \geq a \cdot \mathbb{E}[X]} \Pr(X=i)$$

$$\leq \sum_{i \geq a \cdot \mathbb{E}[X]} i \cdot \Pr(X=i)$$

$$\leq \mathbb{E}[X]$$

(non-neg)

Skip
+
Unroll

Linear Scaling, only needs non-negative

Corollary: Chebyshev's Inequality

$$\Pr(|X - \mathbb{E}[X]| \geq a) \leq \frac{\text{Var}(X)}{a^2}$$

Let $Y = (X - \mathbb{E}[X])^2$ so $\mathbb{E}[Y] = \text{Var}(X)$

$\sqrt{Y} \geq a$
 $\Pr(|X - \mathbb{E}[X]| \geq a) = \Pr(Y \geq a^2)$

$$= \Pr\left(Y \geq \frac{a^2}{\mathbb{E}[Y]} \cdot \mathbb{E}[Y]\right)$$

$$\leq \frac{\mathbb{E}[Y]}{a^2} \quad (\text{Markov})$$

$$= \frac{\text{Var}(X)}{a^2}$$

→ skip + come back if time

Quadratic scaling, needs low variance

Simple Chernoff

Let $X_1, X_2, \dots, X_n$ be independent RVs s.t. $X_i = \begin{cases} 1 & \text{w/ prob. } p \\ 0 & \text{o/w} \end{cases}$

Let $X := \sum_i X_i$, $\mu := \mathbb{E}[X]$

Then $\forall \delta \geq 0$

$$\Pr(X \geq (1+\delta) \cdot \mu) \leq \exp(-\mu \delta^2 / (2+\delta))$$

And $\forall \delta \in (0, 1)$

$$\Pr(X \leq (1-\delta) \cdot \mu) \leq \exp(-\delta^2 \mu / 2)$$

Exponential scaling, needs indicator $\sum_i$

$M \cdot \delta \approx C \cdot \log n \rightarrow \Pr \leq \frac{1}{n^c} \rightarrow \text{"with high probability"}$

Solving Problems w/ Concentration Framework

1) Fix bin i .

$$X_i = \sum_j X_{ij} \quad \left\{ \begin{array}{l} 1 \text{ w/ pr } \frac{1}{n} \\ 0 \text{ o/w} \end{array} \right. \quad \text{independent}$$

(a)

$$\mu = \mathbb{E}[X_i] = 1$$

(b)

$$\begin{aligned} \Pr(X_i \geq \mu(1 + \underbrace{2}_{\delta} \log n)) &\leq \exp(-4 \log^2 n / (2 + \log n)) \\ &\leq \exp(-4 \log^2 n / 2 \cdot \log n) \\ &= n^{-2} \end{aligned}$$

(c)

$$\begin{aligned} \Pr(X_1 \geq 1 + 2 \log n \cup X_2 \geq 1 + 2 \log n \cup \dots) &\leq \sum_i \Pr(X_i \geq 1 + 2 \log n) \\ &\leq \frac{1}{n} \end{aligned}$$

$\mathbb{E}[X_i] = O(1)$ so $X_i \leq \log n$ w.h.p so all $X_i \leq \log n$

2) (a) Let $m = \log n \cdot \ln n$ and fix i

$$\mu = \mathbb{E}[X_i] = \log n$$

(b)

$$\begin{aligned} \Pr(X_i \leq (1 - \underbrace{\frac{1}{2}}_{\delta}) \cdot \log n) &\leq \exp(-\frac{1}{4} \log^2 n / 2) \\ &= n^{-2} \end{aligned}$$

(c)

$$\Pr(X_1 \leq \frac{1}{2} \log n \cup X_2 \leq \frac{1}{2} \log n \cup \dots) \leq \frac{1}{n}$$

$\mathbb{E}[X_i] = \Theta(\log n)$

so $X_i = \Theta(\log n)$ w.h.p

so all $X_i = \Theta(\log n)$ w.h.p

Full Chernoff (Upper Tail)

Let $X_1, X_2, \dots, X_n$ be independent RVs s.t. $X_i = \begin{cases} 1 & \text{w/ prob. } p_i \\ 0 & \text{o/w} \end{cases}$

Let $X := \sum_i X_i$, $\mu := \mathbb{E}[X]$

Then $\forall \delta \geq 0$, $\Pr(X \geq (1+\delta) \cdot \mu) \leq \left(\frac{\exp(\delta)}{(1+\delta)^{1+\delta}} \right)^\mu$

Let $s = \ln(1+\delta)$ and $a = (1+\delta)\mu$

$$\Pr(X \geq a) = \Pr(e^{sX} \geq e^{sa}) = \Pr\left(e^{sX} \geq \frac{e^{sX}}{\mathbb{E}[e^{sX}]}\right) \leq \frac{\mathbb{E}[e^{sX}]}{e^{sa}} = \frac{\mathbb{E}\left[\prod_i e^{sX_i}\right]}{e^{sa}}$$

$\uparrow$ Markov $\uparrow$ def.

$$= \frac{\prod_i \mathbb{E}[e^{sX_i}]}{e^{sa}} = \frac{\prod_i p_i e^s + (1-p_i)}{e^{sa}}$$

$\uparrow$ independence $e^{sX_i} = \begin{cases} e^s & \text{w/ prob. } p_i \\ 1 & \text{w/ prob. } 1-p_i \end{cases}$

$$= \prod_i \frac{1 + p_i(e^s - 1)}{e^{sa}} \leq \prod_i \frac{\exp(p_i(e^s - 1))}{e^{sa}}$$

$\uparrow 1+x \leq \exp(x)$

$$= \frac{\exp(\sum_i p_i(e^s - 1))}{e^{sa}} = \frac{\exp(\mu(e^s - 1))}{e^{sa}}$$

$$= \left(\frac{\exp(\delta)}{(1+\delta)^{1+\delta}} \right)^\mu$$

$\uparrow$

$$s = \ln(1+\delta)$$

$$a = (1+\delta)\mu$$

Proof of Simple Chernoff via Full Chernoff

$$\text{Have } \frac{2\delta}{2+\delta} \leq \log(1+\delta) \quad \forall \delta \geq 0 \quad \xrightarrow{\delta = \delta-1} \quad 2\delta \leq \underbrace{(1+\delta) \cdot \log \delta + 2}_{\text{convex and LHS tangent at } \delta=1 \text{ to RHS}} \quad \forall \delta \geq 1$$

$$\begin{aligned} \text{So } \left(\frac{\exp(\delta)}{(1+\delta)^{1+\delta}} \right)^M &= \left(\frac{\exp(\delta)}{\exp((1+\delta) \cdot \log(1+\delta))} \right)^M \leq \left(\frac{\exp(\delta)}{\exp\left((1+\delta) \cdot \frac{2\delta}{2+\delta}\right)} \right)^M \\ &= \left(\exp\left(\delta - (1+\delta) \frac{2\delta}{2+\delta}\right) \right)^M = \exp\left(-\frac{\delta^2 M}{2+\delta}\right) \end{aligned}$$

Hoeffding's Inequality

Let $X_1, X_2, \dots, X_n$ be independent RVs s.t. $X_i \in [a_i, b_i]$

Let $X := \sum_i X_i$, $\mu := \mathbb{E}[X]$

$$\Pr(|X - \mu| \geq t) \leq 2 \cdot \exp\left(-\frac{2t^2}{\sum_i (b_i - a_i)^2}\right)$$

Like Chernoff but additive/doesn't assume $X_i \in \{0, 1\}$